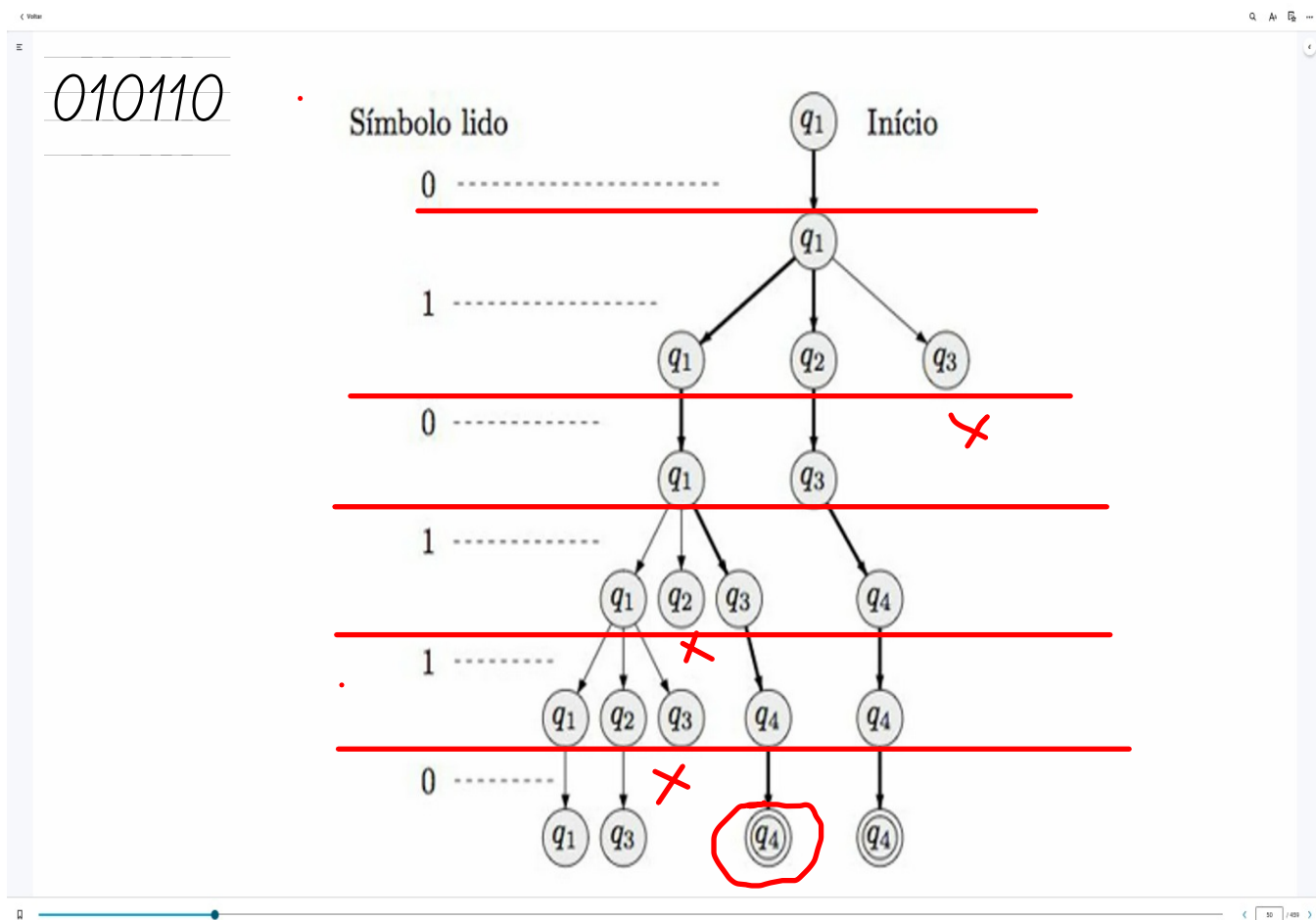


FIGURA 1.27
O autômato finito não-determinístico N_1 .



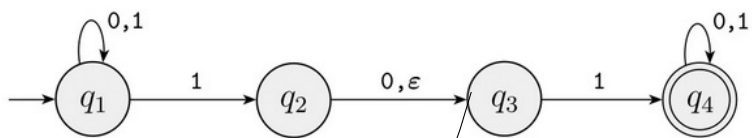


FIGURA 1.27
O autômato finito não-determinístico N_1 .

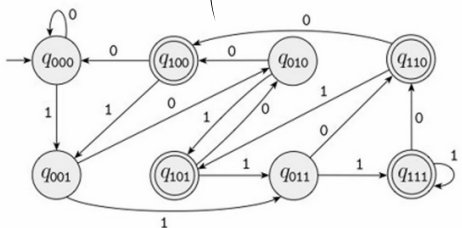
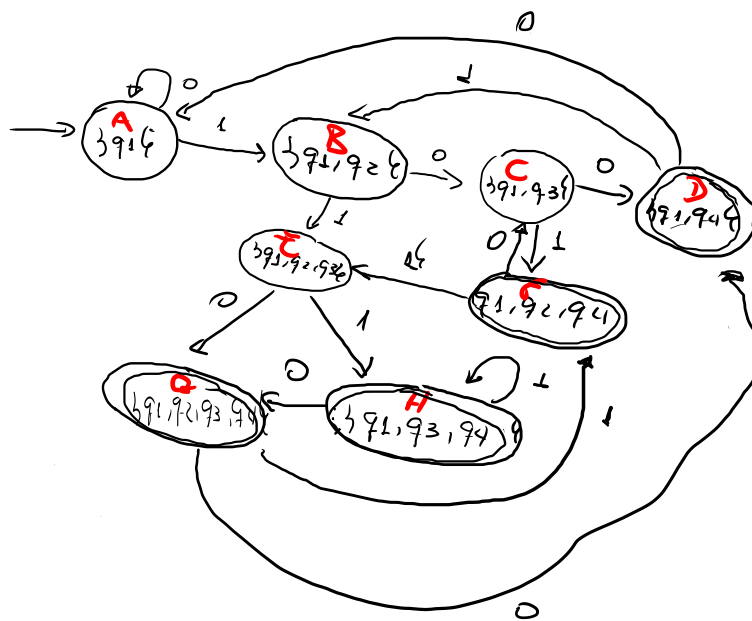


FIGURA 1.32
Um AFD que reconhece A .



FIGURA 1.31
O AFN N_2 que reconhece A .



$E()$
 ϵ -move

34

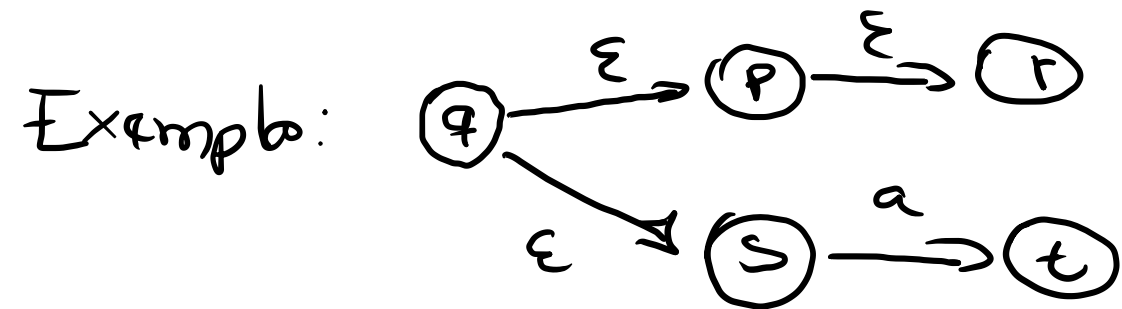
2^n

$$\delta(q_1, 0) = \{q_1\}$$

$$\delta(q_1, 1) = \{q_1, q_2\}$$

FUNÇÃO $E()$

$E(q)$ é o conjunto de todos os estados alcançáveis com ϵ , incluindo q .



$$E(q) = \{q, p, r, s\}$$

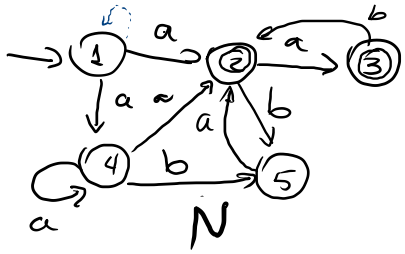
$$E(\{p_1, p_2, \dots, p_m\}) = E(p_1) \cup E(p_2) \dots E(p_m)$$

No exemplo acima:

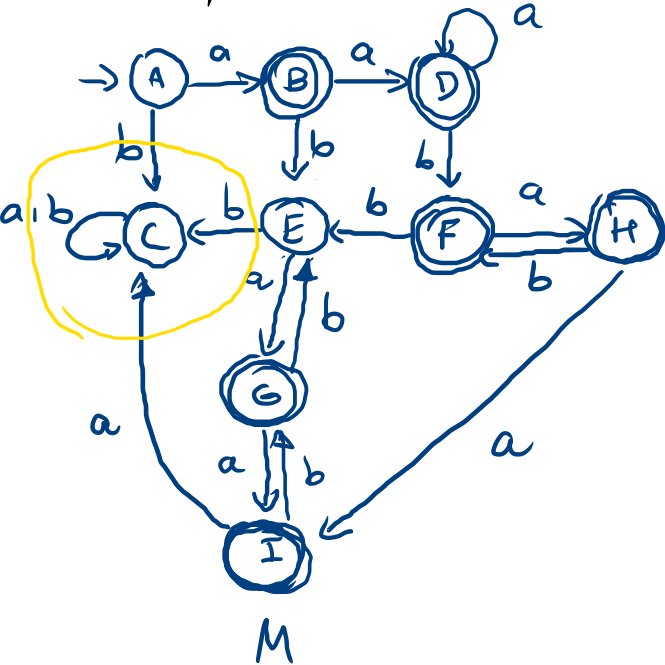
$$\begin{aligned} E(p, s) &= E(p) \cup E(s) \\ &= \{p, r\} \cup \{s\} = \{p, r, s\} \end{aligned}$$

Seja o seguinte AFN representado abaixo

$$\delta(p_1, p_2, \dots, p_n, a) = \delta(p_1, a) \cup \delta(p_2, a) \dots$$



$L(N) \equiv L(M)$



$$E(L) = \{1\} \oplus E(\delta(\{1\}, a)) = E(\{2, 4\}) = \{2, 4\} \oplus$$

$$E(\delta(\{1\}, b)) = E(\{4\}) = \{4\} \oplus$$

$$E(\delta(\{2, 4\}, a)) = E(\{3, 2, 4\}) = \{2, 3, 4\} \oplus$$

$$E(\delta(\{2, 4\}, b)) = E(\{5\}) = \{5\} \oplus$$

A função E p/ transições sem ϵ é iterativa.

$$E(\delta(\{4\}, a)) = \{4\} \quad E(\delta(\{2, 3, 4\}, a)) = \{2, 3, 4\}$$

$$E(\delta(\{4\}, b)) = \{4\} \quad E(\delta(\{2, 3, 4\}, b)) = \{2, 5\} \oplus$$

$$E(\delta(\{5\}, a)) = \{2\} \oplus \quad E(\delta(\{2, 5\}, a)) = \{2, 3\} \oplus$$

$$E(\delta(\{5\}, b)) = \{4\} \quad E(\delta(\{2, 5\}, b)) = \{5\}$$

$$E(\delta(\{2\}, a)) = \{3\} \oplus \quad E(\delta(\{2, 3\}, a)) = \{3\}$$

$$E(\delta(\{2\}, b)) = \{5\} \quad E(\delta(\{2, 3\}, b)) = \{2, 5\}$$

$$E(\delta(\{3\}, a)) = \{4\}$$

$$E(\delta(\{3\}, b)) = \{2\}$$

Os conjuntos finais são todos que contêm antigos finais.

A = {1}	<
B = {2, 4}	<
C = {4}	<
D = {2, 3, 4}	<
E = {5}	<
F = {2, 5}	<
G = {2}	<
H = {2, 3}	<
I = {3}	<